

Homework #3

(Search)

Out: 1/24/11 — Due: 2/1/11 (at noon)

How to complete this HW: First copy this file; then type your answers in the file immediately below each question; start each question on a separate page, write your name on every page. Finally, print this file and return it in the drawer marked CS121 of a file cabinet located in the entryway of the Gates building next to office Gates 182 (see red arrow in <http://ai.stanford.edu/~latombe/cs121/2011/map.jpg>) no later than Tuesday 2/1 at noon.

Your name:

Your email address:

Note on Honor Code: You must NOT look at previously published solutions of any of these problems in preparing your answers. You may discuss these problems with other students in the class (in fact, you are encouraged to do so) and/or look into other documents (books, web sites), with the exception of published solutions, without taking any written or electronic notes. If you have discussed any of the problems with other students, indicate their name(s) here:

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Any intentional transgression of these rules will be considered an honor code violation.

General information: Justify your answers, but keep explanations short and to the point. Excessive verbosity will be penalized. If you have any doubt on how to interpret a question, tell us in advance, so that we can help you understand the question, or tell us how you understand it in your returned solution.

Grading:

Problem#	Max. grade	Your grade
I	25	
II	25	
III	25	
IV	25	
Total	100	

Name

I. Business travel in California (25 points)

You are given a road map of California in the form of a graph in which each node is a city and each edge represents a road between two cities. The graph is such that every two nodes are connected by a path made of one or several consecutive edges. Each edge (i,j) is labeled by the length $l(i,j)$ of the road between cities i and j . The two pairs (i,j) and (j,i) designate the same edge; in other words, edges are non-directed. Any two cities connected by an edge in the map are said to be *adjacent*.

Two business persons A and B start from two different cities of the map, respectively a and b . They want to meet in a city of the map (any one) to discuss a new business project. To achieve this goal, they move in successive steps. At each step, both A and B start moving at the same time and each must move from some city i to an adjacent city j . Neither A nor B can stay in the same city.

We assume that both A and B move at unit velocity. So, the amount of time to move from city i to neighboring city j is equal to the length $l(i,j)$ of the road between the two cities. Since A and B do not usually travel the same distance, at the end of each step, they may not reach their respective new cities at the same time. The person that arrives first to his/her new city must wait until the other arrives to his/her new city (each one calls the other on his/her cell phone when he/she arrives to a new city) before the next step can begin.

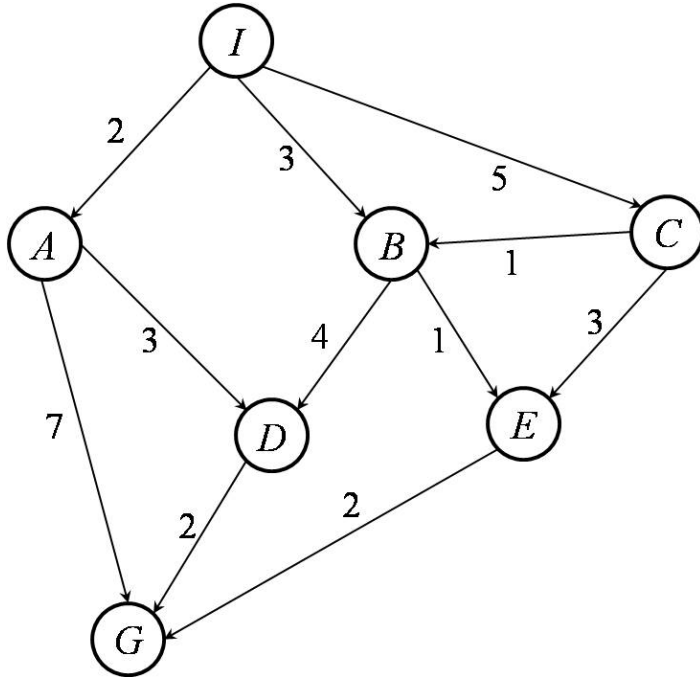
Under the above conditions, the two business persons want to meet as quickly as possible. Recall that their goal is to meet in a city of the map (any one), but not anywhere on a road.

1. Formulate this problem as a state-space search problem:
 - a) What is the state space?
 - b) What are the initial and the goal states?
 - c) What is the successor function?
 - d) What is the cost function for an arc?
2. Let $D(i,j)$ be the straight-line distance between any two cities i and j in the map. Assume that D is known for each pair of cities. Using D , define the most informed monotone admissible heuristic function h for the search problem you have defined in question 1.
3. Build a road map with $n > 1$ cities (where n is arbitrary) satisfying the description of the first paragraph of this problem such that no solution exists for certain pairs of initial cities a and b . Note that the first paragraph says: "The graph is such that every two nodes are connected by a path made of one or several consecutive edges."

Name

II. Best-First Search (25 points)

Consider the search problem defined by the following search graph in which I is the initial state and G is the only goal state.



Node	h_0	h_1	h_2	h_3
I	1	6	9	6
A	0	5	8	5
B	1	3	7	3
C	0	4	5	5
D	1	2	3	2
E	1	2	3	2
G	0	0	0	0

The cost of each arc is indicated beside each arc. The table of the right gives the values of 4 heuristic functions h_i , $i = 0, 1, 2$, and 3.

- Which of the heuristics shown in the table are admissible? Which ones are consistent? No need to justify your answers.
- Let us assume that use best-first search with the evaluation function $f_i(N) = g(N) + h_i(N)$ and that the search algorithm tests if a node is a goal node when it expands this node. What sequence of states (connecting I to G) would be returned by the algorithm using each of the heuristics? Assume that when several nodes in the fringe share the same minimum value of f_i , then the algorithm expands the nodes in alphabetic order of their states.

Name

III. Robot on a grid (25 points)

A robot moves from vertices to vertices in the unbounded regular 2D grid shown in Figure 1. The initial position of the robot is the vertex $(0,0)$ marked with a black dot. At each step the robot can move from its current position by a unit increment up, down, left, or right. The goal is for the robot to move to some given grid vertex (x,y) .

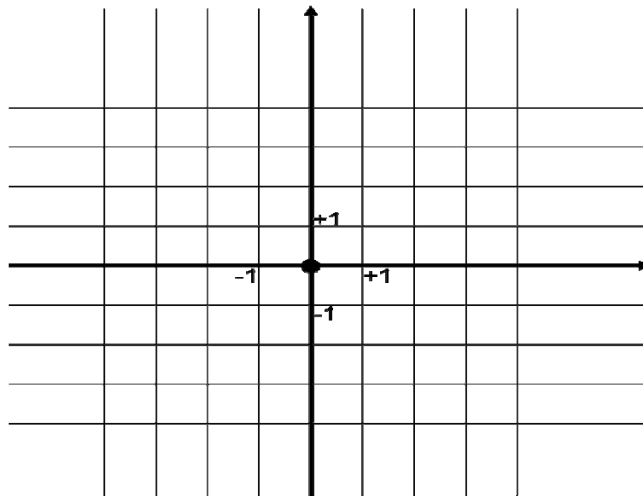


Figure 1: A portion of the grid used in Problem II

- How many distinct positions can the robot reach in $k > 0$ steps that it could not reach in less than k steps?
(i) 4^k (ii) $4k$ (iii) $4k^2$
- Let $n = |x| + |y|$. How many nodes, at most, does breadth-first search *without* revisited-state checking *expand* before terminating? [You *must* explain your answer.]
(i) $((4^{n+1} - 1)/3) - 1$ (ii) $4n^2 - 1$ (iii) $2n(n+1) - 1$
- Let $n = |x| + |y|$. How many nodes, at most, does breadth-first search *with* revisited-state checking *expand* before terminating? [You *must* explain your answer.]
(i) $((4^{n+1} - 1)/3)$ (ii) $4n^2$ (iii) $2n(n+1)$
- Is $h(N) = |u-x| + |v-y|$ a consistent heuristic, where (u,v) is the state of node N ? Why?
- Does A* with revisited-state checking using h expands $O(|x| + |y|)$ nodes before terminating?
- Does the heuristic h defined above remain consistent if some links in the grid are removed?
- Does h remain consistent if some links are added?

Name

IV. Approximately optimal search (25 points)

The two objectives of (1) finding a solution as quickly as possible and (2) finding an optimal solution are often conflicting. In some problems, one may design two heuristic functions h_A and h_N , such that h_A is admissible and h_N is not admissible, with h_N resulting in much faster search most of the time. Then, one may try to take advantage of both functions.

Let $f(N) = g(N) + h_A(N)$, where $g(N)$ is the cost of a path from the initial node to node N . Let A_ϵ^* search be defined as follows. Like A^* , it runs the SEARCH#2 algorithm, hence it terminates when the node selected for expansion is a goal node. Also, like A^* (since h_A is only known to be admissible), A_ϵ^* does not try to detect if states are re-visited. However, unlike A^* , at each expansion step, A_ϵ^* expands a node N' (any one) such that $f(N') \leq (1+\epsilon) \times \min_{N \in \text{FRINGE}} \{f(N)\}$, where ϵ is any strictly positive number. In other words, if N is the node in the fringe that has the smallest value of f , then A_ϵ^* expands a node N' such that $f(N) \leq f(N') \leq (1+\epsilon) \times f(N)$.

1. Prove that A_ϵ^* terminates whenever a solution exists.
2. What can you say about the cost of the solution returned by A_ϵ^* ? [*Hint: Prove that the cost of the path found by is only slightly greater than the cost of the optimal path.*]
3. Explain briefly how A_ϵ^* can use the second heuristic function h_N to reduce the time of the search. What tradeoff is being made in choosing ϵ ? [*Hint: What are the advantages and drawbacks of increasing the value of ϵ ?*]